

Methods of proof, Lec 7.

3. Proof by contradiction of a proposition p

a) Find a second proposition, r , that is implied from p . $p \rightarrow r$

Typically r is a definition associated w/ p .

e.g. Proof $\sqrt{2}$ is irrational

required defn $r = \text{rational number } r = l/m$

where $l, m \in \mathbb{Z}$ w/ no common factors

b) show that $\neg p \rightarrow (r \wedge \neg r)$ a contradiction.

$$\text{Thus } \neg p = F \text{ and } \neg \neg p = p = T //$$

(an indirect proof).

3. Proof by contradiction of an implication $p \rightarrow q$

Foundations

$$\neg(p \rightarrow q) \equiv p \wedge \neg q \quad \begin{matrix} p \rightarrow q \\ \text{(only false if } p=T \text{ and } q=F) \end{matrix}$$

$$\neg(p \rightarrow q) \equiv \neg(\neg q \rightarrow \neg p) \equiv \neg q \wedge p \equiv p \wedge \neg q$$

\uparrow
commutativity.

Need to show if $p \rightarrow q$ is False, this leads to a contradiction.

Proof: show $\neg(p \rightarrow q)$ leads to $(p \wedge \neg p)$ or $(q \wedge \neg q)$

$$\begin{array}{l} p \\ \neg q \\ \hline \end{array} \quad \begin{array}{l} \text{assert } p \\ \text{assert } \neg q \end{array}$$
$$\therefore (p \wedge \neg p) \vee (q \wedge \neg q)$$

"Domination law" $p \vee T \equiv T$, show either $(p \wedge \neg p)$ or $(q \wedge \neg q)$

$$\text{since } \neg(p \rightarrow q) \rightarrow (p \wedge \neg p) \vee (q \wedge \neg q)$$

$$\text{therefore } \neg(p \rightarrow q) = F \text{ and } p \rightarrow q = T$$

Example proof: by contradiction of proposition

$$p = 4 \text{ is an odd number}$$

$$\text{If } p \text{ is an odd number, } \rightarrow r = n = 2k + 1$$

$$\text{for some } k \in \mathbb{Z}$$

$$n \text{ is an even number}$$

$$\rightarrow q = n = 2k, \text{ for some } k \in \mathbb{Z}$$

Argument form:

$$\neg p \rightarrow (r \wedge \neg r)$$

Proof:

$$4 = 2j \quad (\text{assert } \neg p)$$

No integer can be both odd & even.

$$\therefore r \wedge \neg r$$

$$4 = 2j \text{ and } 4 = 2k + 1$$

Example: Proof by contradiction of $p \rightarrow q$

Proof: $\forall x, y \in \mathbb{R} \left(\underbrace{(x+y) \geq 2}_{P(x,y)} \rightarrow \underbrace{(x \geq 1 \vee y \geq 1)}_{Q(x,y)} \right)$

Recall argument form

$$\begin{array}{c} p \\ \neg q \\ \hline \therefore (p \wedge \neg p) \vee (q \wedge \neg q) \end{array}$$

$$x+y \geq 2 \quad (\text{assert } P \text{ is true})$$

$$\neg (x \geq 1 \vee y \geq 1) \quad (\text{assert } \neg q)$$

$$(\neg (x \geq 1) \wedge \neg (y \geq 1)) \quad \text{DeMorgan's}$$

$$(x < 1) \wedge (y < 1) \quad \text{arithmetic.}$$

$$\therefore p \wedge \neg p$$

aside: $\left\{ \begin{array}{l} p = (x \geq 1) \vee (y \geq 1) \\ \neg q = (x < 1) \wedge (y < 1) \end{array} \right\}$

$$\begin{array}{l} p = x+y \geq 2 \\ \text{if } (x < 1) \wedge (y < 1) \\ \hline \therefore \neg p \end{array}$$

$$\text{Since } \neg(p \rightarrow q) \equiv p \wedge \neg q = F$$

$$\begin{array}{l} \text{since } x+y < 2 \\ \text{if } (x < 1) \wedge (y < 1) \end{array}$$

Then $p \rightarrow q$ is true.

4. Equivalence proof (if and only if)

$$p \leftrightarrow q$$

Need to show $p \rightarrow q$ and $q \rightarrow p$.

Example: If $n \in \mathbb{Z}^+$ then n is odd ~~only~~ if and only if n^2 is odd.

1) $p \rightarrow q$ direct proof easy.

$$n = 2k+1 \quad (\text{assert } p)$$

$$n^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1 = 2j + 1$$

$\therefore n^2$ is odd.

2) $q \rightarrow p$ try direct proof

$$n^2 = 2k+1 \quad (\text{assert } q)$$

$$n = \sqrt{2k+1} \dots ???$$

Try contrapositive $\neg p \rightarrow \neg q$

$$n = 2k \quad (\neg p \text{ assert})$$

$$n^2 = 4k^2 = 2(2k^2) = 2j$$

$\therefore n^2$ is even; $\neg p \rightarrow \neg q \therefore q \rightarrow p$ //

5. Exhaustive proof / proof by cases.

* show that the implication is true for all $x \in U$.

Very often the elements x fall into classes:

e.g. integers \mathbb{Z} ; ^{positive} ~~odd~~, ^{negative} ~~even~~, zero (3 classes)

• real #'s \mathbb{R} ; rational or irrational (2 classes)

Proof that $p \rightarrow q$ holds for each category.

(Challenge is to make sure you cover all the categories).

Example; proof if n is an integer then $n^2 \geq n$.

Prove: $\forall n \in \mathbb{Z} (n^2 \geq n)$

3 classes of integers: - positive
- negative
- zero

• case 1 $n > 0$ then $n^2 \geq n$; $\therefore \frac{n \geq 0}{n^2 \geq n}$

• case 2 $n < 0$ then $n^2 > n$; $\therefore \frac{n < 0}{n^2 > n}$

• case 3 $n = 0$ then $n^2 = 0$ and $n^2 \geq n$. $\therefore \frac{n = 0}{n^2 \geq n}$

Often proof by cases relies on

integers being

• odd or even

• positive, negative, or zero.

Existence proofs

$$\exists x P(x)$$

$$\exists x (P(x) \rightarrow Q(x))$$

6. Constructive existence proofs ; find an x for which $P(x)$.

7. Nonconstructive existence proofs;
show $P(x)$ must hold for some x
but we don't need to specify
the x precisely.

Example: constructive existence proof

Prove there exist an ^{positive} integer n that is the
sum of cubes of positive integers in two ways.

$$\exists n \in \mathbb{Z}^+ \text{ s.t. } (n = i^3 + j^3) \wedge (n = k^3 + l^3) \\ \text{for } (i \neq k \neq l, j \neq k, l.)$$

$$1729 = 10^3 + 9^3 = 12^3 + 1^3 //$$

↑
in other
words
 $i \neq k \neq l \neq j$

Example of

#7: Non constructive existence proof.

* Show there exists a value, but don't have to pinpoint it.

Prove: that there exists irrational x, y .

s.t. x^y is rational.

Recall rational # $r = l/m$, where $l, m \in \mathbb{Z}$ w. no common factors.

Need an instance where x, y are both irrational but x^y is rational.

Try $x = \sqrt{2}$ and $y = \sqrt{2}$

consider $x^y = \sqrt{2}^{\sqrt{2}}$

case 1: if $\sqrt{2}^{\sqrt{2}}$ is rational \rightarrow done w/ existence proof.

case 2: if $\sqrt{2}^{\sqrt{2}}$ is irrational

let $x = \sqrt{2}^{\sqrt{2}}$ and $y = \sqrt{2}$

$$x^y = (\sqrt{2}^{\sqrt{2}})^{\sqrt{2}} = (\sqrt{2})^2 = 2.$$

$(\sqrt{2}^{\sqrt{2}})^{\sqrt{2}}$ is rational