

Lec 14. Proof by induction

Proof technique for proving propositions $P(n)$
about $n \in \mathbb{N}, \mathbb{Z}^+$

Three steps:

- 1) Assume $P(n)$ (all defns & associated axioms)
- 2) Use the defns of $P(n)$ to show
that if $P(n)$ is true, then $P(n+1)$ is true.
Direct proof $P(n) \rightarrow P(n+1)$.
- 3) Show $P(n)$ holds for some smallest
element \mathbb{Z}^+ or \mathbb{N} , call it n_{\min} .

$$\therefore \forall n \geq n_{\min} \quad P(n) = T.$$

Typically useful for:

- summations
- Inequalities
- Division
- Sets
- Algorithms.

Proof by induction for summations (General technique).

Example summations

$$S_n = \sum_{i=1}^n i, \quad S_{\text{geo}}(n) = \sum_{i=0}^n ar^i, \quad \text{etc.}$$

Hypothesis $P(n) := \sum_n^v = f(n)$
raw sum \nearrow simple, closed-form, eqn.

Need to show $P(n) \rightarrow P(n+1)$

i.e. if $\sum_n^v = f(n)$ then $\sum_{n+1}^v = f(n+1)$

Steps:

- Write \sum_{n+1}^v and execute algebra to make it expressed in terms of \sum_n^v .
- Substitute $f(n)$ for \sum_n^v (i.e. assume $P(n)$)
- Use algebra to make r.h.s. in terms $f(n+1)$.

For example. Prove $P(n) \rightarrow P(n+1)$

where $P(n) := \sum_{i=1}^n (2i-1) = n^2$

• write S_{n+1} & show it in terms of S_n

$$\sum_{i=1}^{n+1} (2i-1) = \underbrace{\sum_{i=1}^n (2i-1)}_{P(n)} + (2(n+1)-1)$$

• substitute $f(n)$ for $P(n)$

$$\sum_{i=1}^{n+1} (2i-1) = n^2 + (2(n+1)-1)$$

• Use algebra on r.h.s. to show it is $f(n+1)$

$$\sum_{i=1}^{n+1} (2i-1) = n^2 + 2n + 1 = (n+1)^2 //$$

We have proven inductive step $P(n) \rightarrow P(n+1)$

still need the basis step:

$$P(1) = (1)^2 = 1 \quad \checkmark$$

* Some are easy! well defined methodology just shown.

Proof by induction for inequalities:

For example: $P(n) := n < 2^n$ for all $n \in \mathbb{Z}^+$

- Inductive hypothesis, $P(n)$
- Inductive proof step, $P(n) \rightarrow P(n+1)$
- Bases step $P(n_{\min}) = T$

• Hypothesis $P(n) := n < 2^n, \forall n \in \mathbb{Z}^+$

• Inductive proof step

$P(n): \quad n < 2^n$

Wrote
l.h.s of $P(n+1)$

$$\left. \begin{aligned} n+1 &< (2^n)+1 \\ &< 2 \cdot 2^n \\ &= 2^{n+1} \end{aligned} \right\} \text{by Lemma 1}$$

verify

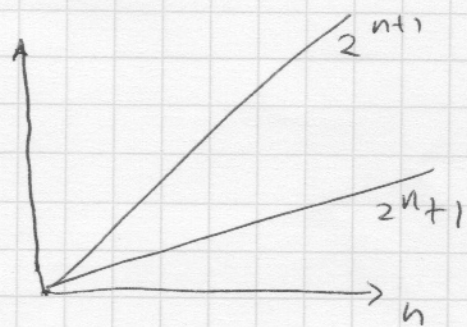
n	$2^n + 1$	$2 \cdot 2^n$
1	3	4
2	5	8
3	9	16
...

$\therefore n+1 < 2^{n+1}$

$(P(n) \rightarrow P(n+1))$

Lemma 1:

For $n > 0, 2^n + 1 < 2^{n+1}$



Proof by induction on sets

For example:

$P(n) :=$ A set with n elements has 2^n subsets.

Want to show (1) $P(n) \rightarrow P(n+1)$; (2) $P(n_{\min}) = T$

1) $P(n) \rightarrow P(n+1)$

Assume $P(n)$ a set of size n has 2^n subsets

Let $|S| = n$

Let $|B| = n+1$ and $B = S \cup \{b\}$ where $b \notin S$

For every subset of S , can make a ^{new} subset w/ $\{b\}$ included.

$$\begin{aligned} \# \text{ subsets of } B &= 2 \# \text{ of subsets of } S \\ &= 2 \cdot 2^n = 2^{n+1} \end{aligned}$$

$P(n) :=$ # of subsets of S is 2^n .

$\therefore P(n) \rightarrow P(n+1)$

2) Bases step.

$P(n)$ holds for some n_{\min}

$P(n=0)$; $|\{\emptyset\}| = 0$, # of subsets $\{\emptyset\} = 2^0 = 1$.

$P(n=0)$ is true (that a set with zero elements has 2^0 subsets).

More proofs for sets:

Hypothesis:

$$P(n) := (A_1 \cup A_2 \cup \dots \cup A_n) \cap B = (A_1 \cap B) \cup (A_2 \cap B) \cup \dots \cup (A_n \cap B)$$

Basis step: try $n=1$.

$$P(1) : A_1 \cap B = A_1 \cap B //$$

Inductive step: Prove $P(n) \rightarrow P(n+1)$

Use distributive law:

$$(X \cup Y) \cap Z = (X \cap Z) \cup (Y \cap Z)$$

Show $P(n) \rightarrow P(n+1)$

$$P(n+1) = \underbrace{(A_1 \cup A_2 \cup \dots \cup A_n)}_X \cup \underbrace{A_{n+1}}_Y \cap \underbrace{B}_Z$$

$\stackrel{=}{\nearrow}$
dist law

$$\underbrace{(A_1 \cup A_2 \cup \dots \cup A_n) \cap B}_{P(n)} \cup (A_{n+1} \cap B)$$

$$= (A_1 \cap B) \cup (A_2 \cap B) \cup \dots \cup (A_n \cap B) \cup (A_{n+1} \cap B)$$

$$= P(n+1) //$$

Recap proof by induction

- Hypothesis $P(n)$ for $n \in \mathbb{N}, \mathbb{Z}^+$
- Inductive step $P(n) \rightarrow P(n+1)$
(using def'n of $P(n)$ and algebra)
- Basis step $P(n) = T$ for some n_{\min}

$$\therefore \forall_{n \geq n_{\min}} P(n)$$

Game plan for $P(n) \rightarrow P(n+1)$:

- Write the l.h.s. of $P(n+1)$, express it in terms of $P(n)$
- Substitute the expression for $P(n)$ with the hypothesis
- Show this results in appropriate formula for r.h.s. of $P(n+1)$

New topic : Recursion and iteration

↖
build from top down.

↗
build up from bottom

Recursion: a general term for defining an object in terms of itself.

- An inductive proof establishes truth of $P(k+1)$ due to truth of $P(k)$.
- A recursive defn of a func, predicate, set, etc. defines larger elements in terms of smaller elements.

For instance consider the geometric sequence

$$\{a_0, a_1, a_2, \dots, a_n\} \text{ where } a_n = ar^n$$

(for a, r const)

A recursive defn of geometric series:

$$\boxed{a_n = (a_{n-1})r}$$

For instance; arithmetic series:

$$\{a_0, a_1, \dots, a_n\} \text{ where}$$

$$a_n = a_0 + nd$$

for a_0, d constants.

Recursive defn.

$$\boxed{a_n = a_{n-1} + d}$$

A recursive func $f(n)$ is defined in terms of smaller values of the func.

e.g. $f(0), f(1), f(2) \dots f(k)$ given

For $n > k$ there is a rule for writing

$f(n)$ in $f(0), f(1), f(2) \dots, f(n-1)$.

Example recursive func's:

Explicit: Factorial func' $f(n) = n \cdot (n-1) \cdot (n-2) \dots 2 \cdot 1$

Recursive factorial func's:

$$f(0) = 1 \quad (\text{Base case})$$

$$\forall n > 0, f(n) = n \cdot f(n-1) \quad \text{Recursive defn.}$$

Explicit func' $f(n) = a^n$

Recursive defn: $f(0) = 1$

$$\forall n > 0, f(n) = a \cdot f(n-1)$$